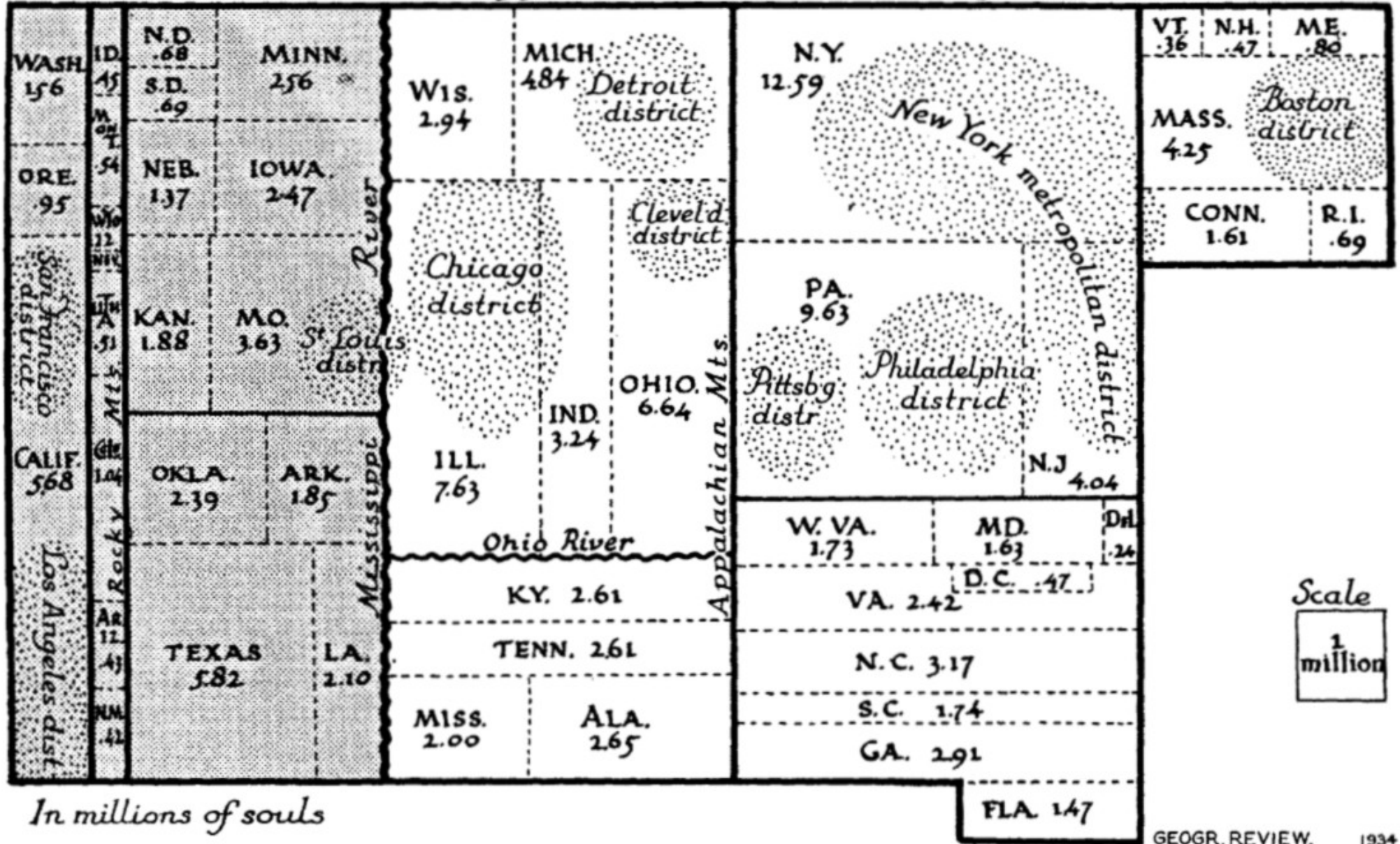


Physics in Unexpected Places

Mark Newman
University of Michigan
and
Santa Fe Institute

POPULATION

1930 census. U.S. total 123.6 million



The diffusion cartogram

We need a process that moves population away from high-density areas into low-density ones until everything is uniform.

$$\mathbf{J} = \mathbf{v}(\mathbf{r}, t) \rho(\mathbf{r}, t) \quad \text{and} \quad \mathbf{J} = -\nabla \rho,$$

The diffusing population is conserved locally:

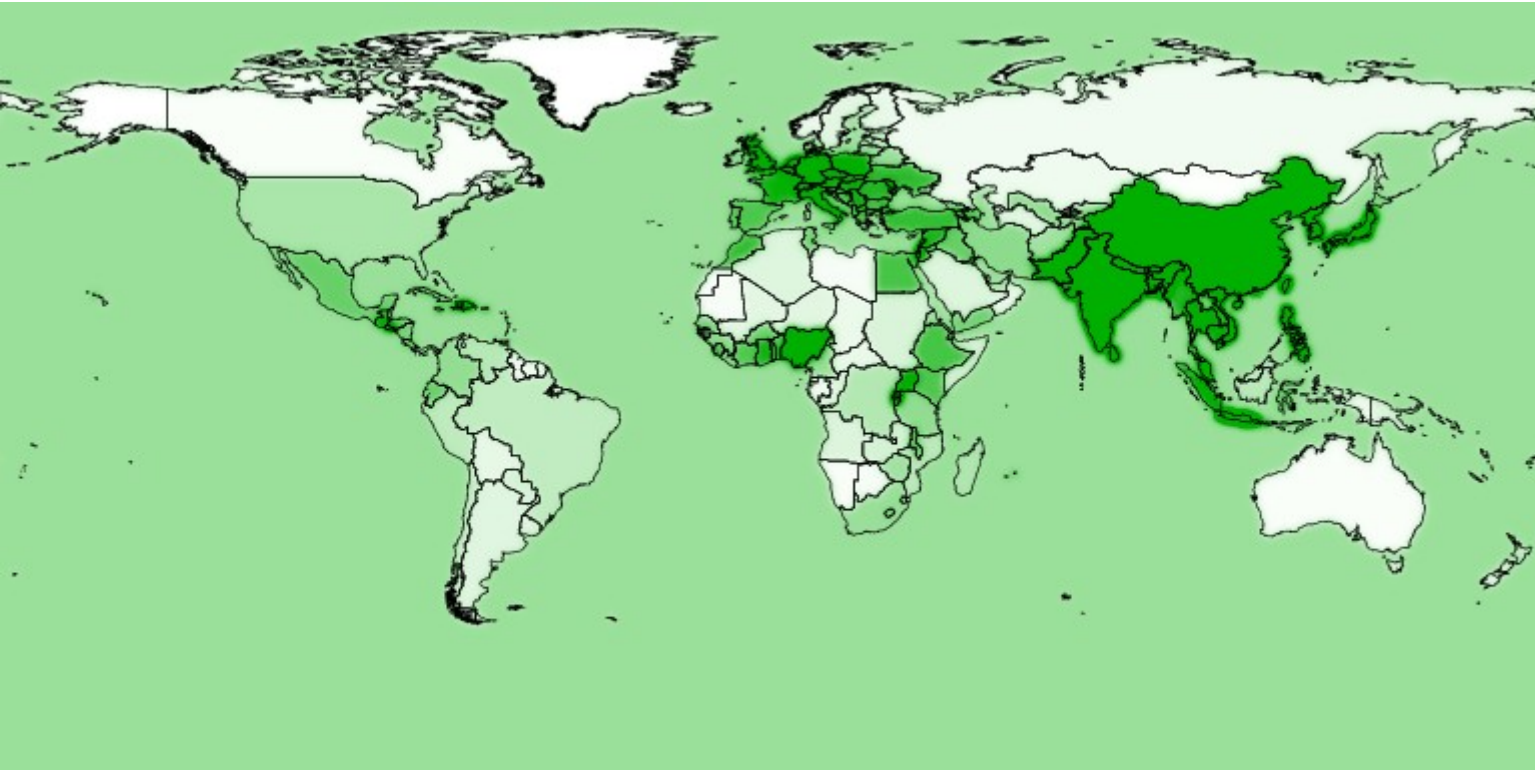
$$\nabla \cdot \mathbf{J} + \frac{\partial \rho}{\partial t} = 0.$$

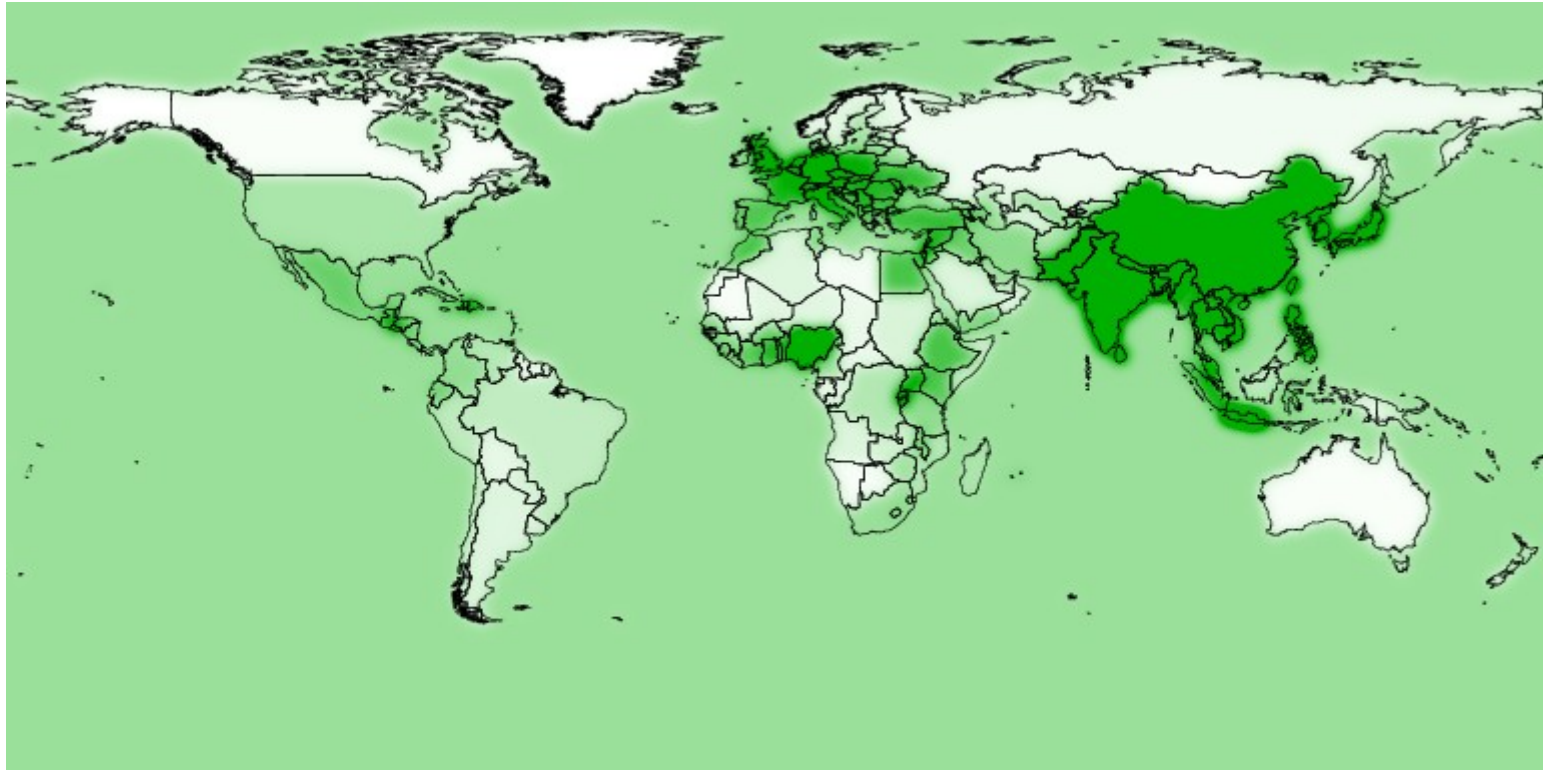
Hence

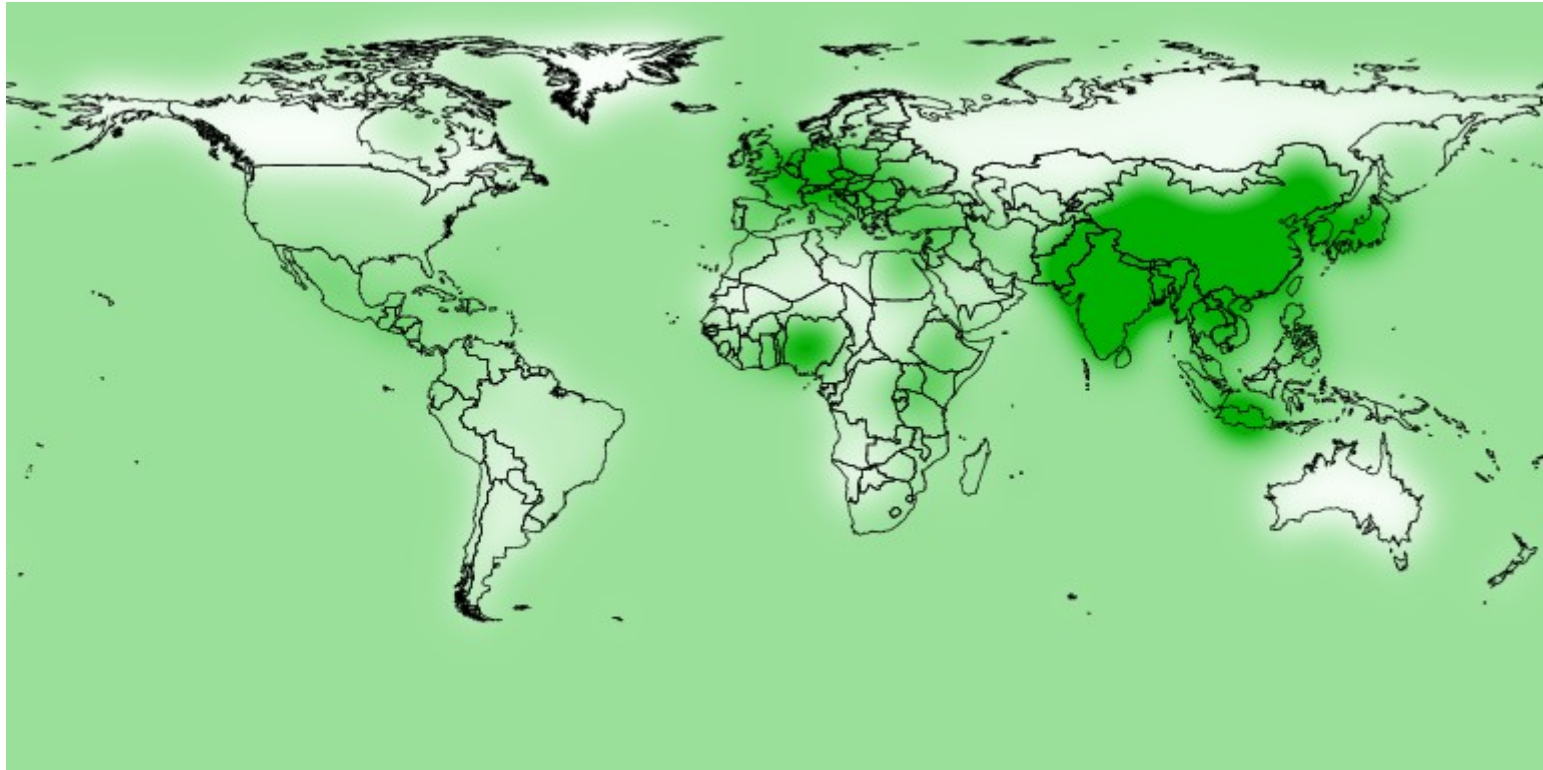
$$\nabla^2 \rho - \frac{\partial \rho}{\partial t} = 0 \quad \text{and} \quad \mathbf{v}(\mathbf{r}, t) = -\frac{\nabla \rho}{\rho}.$$

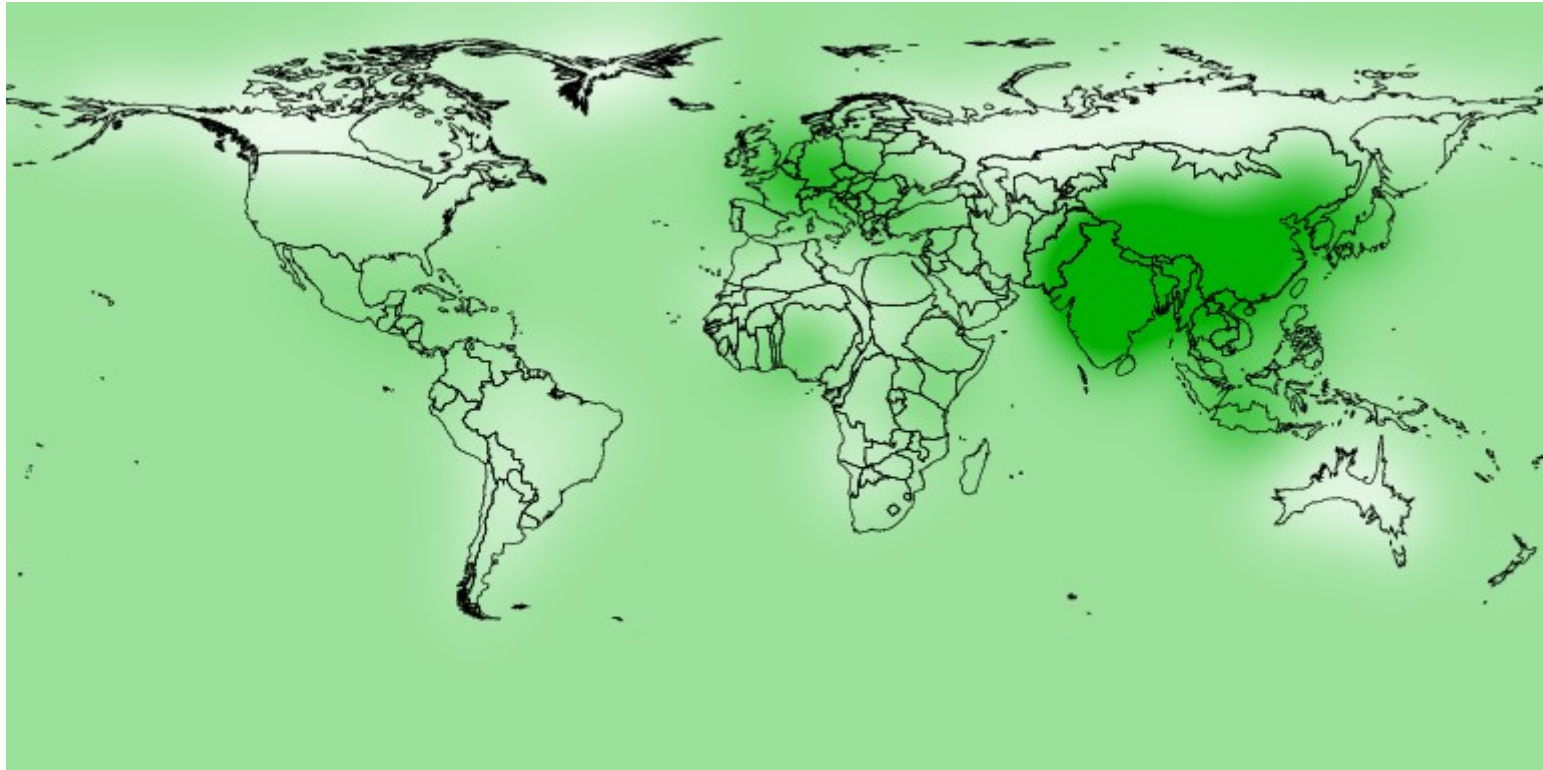


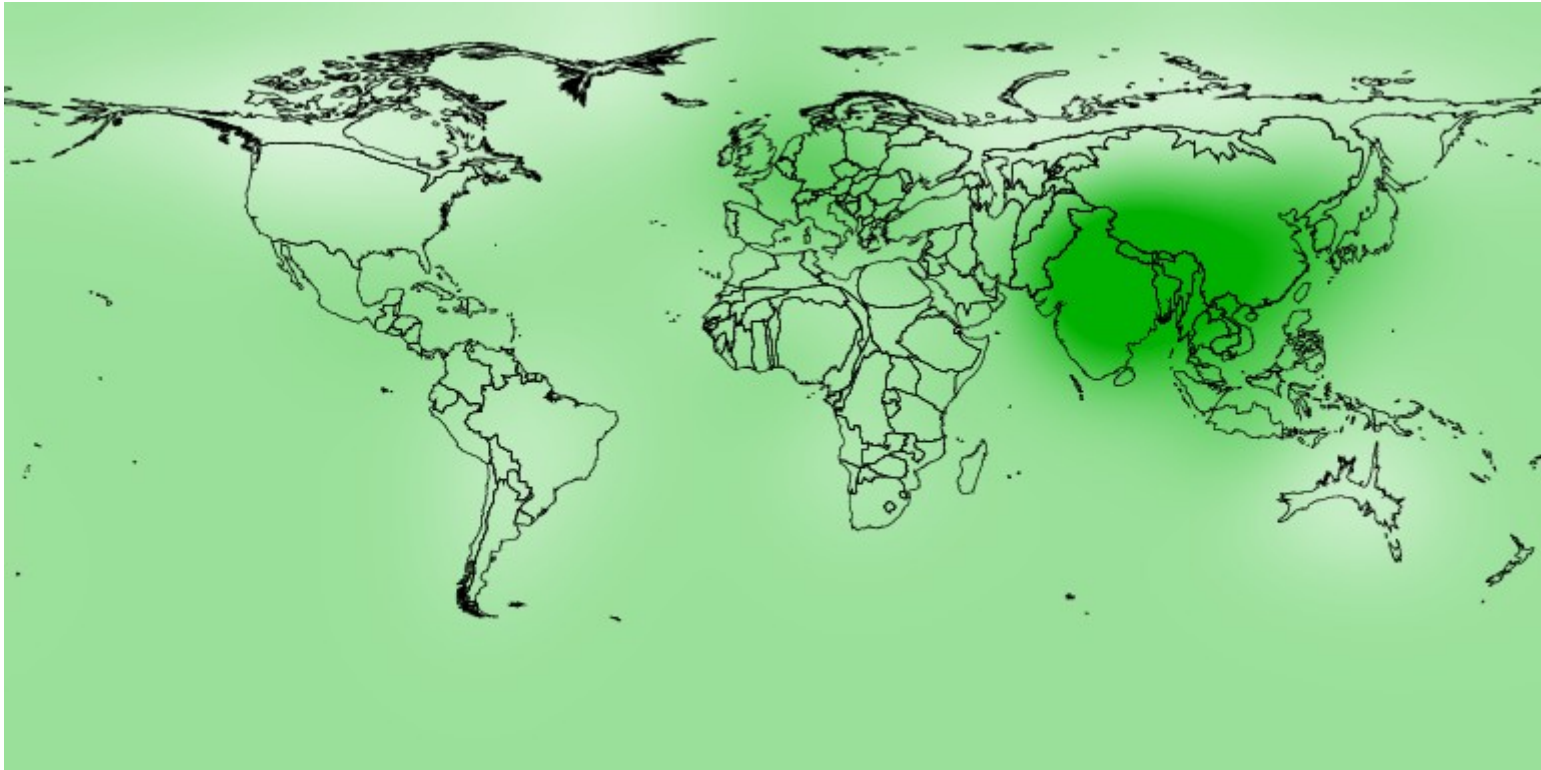


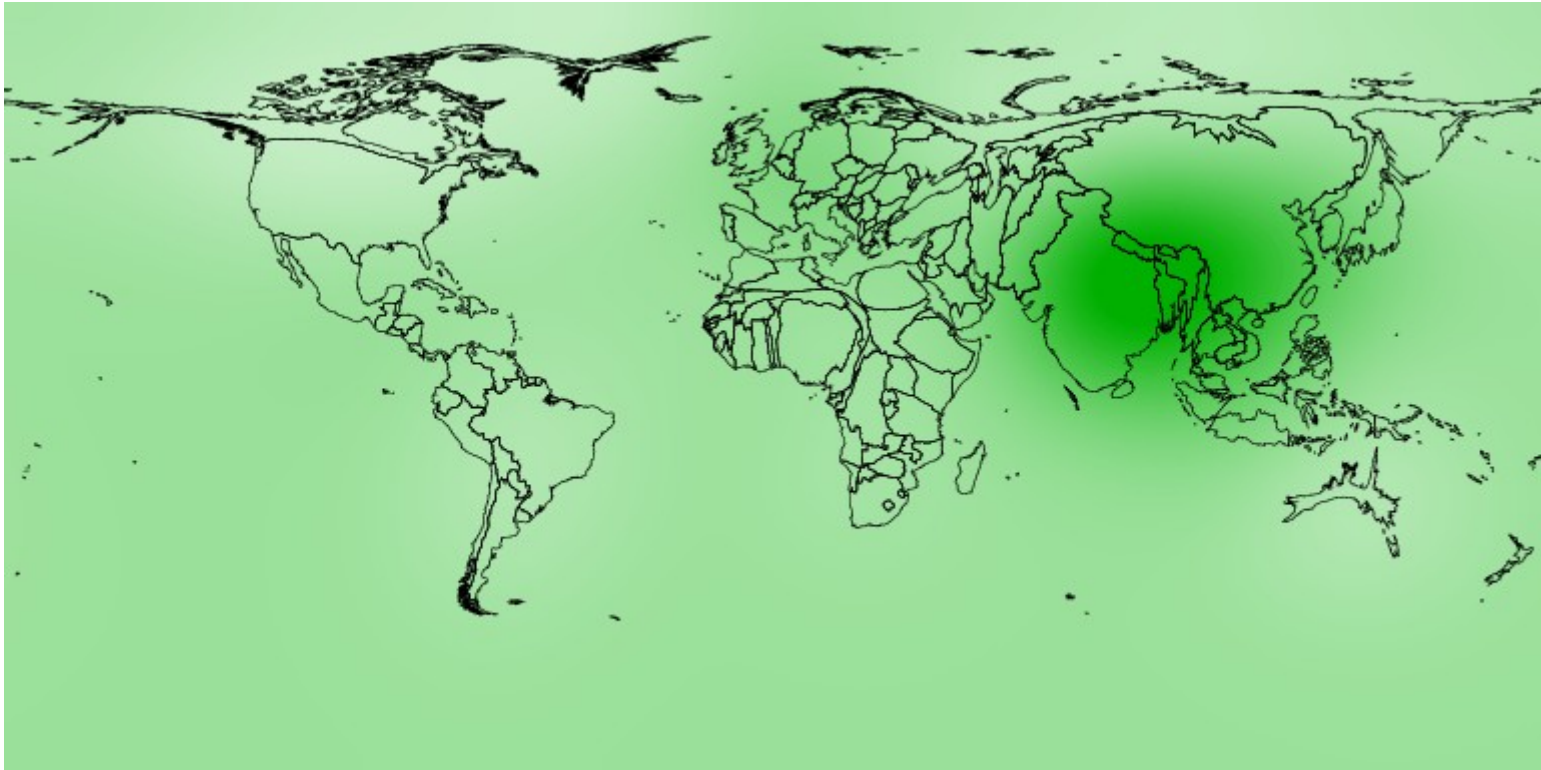


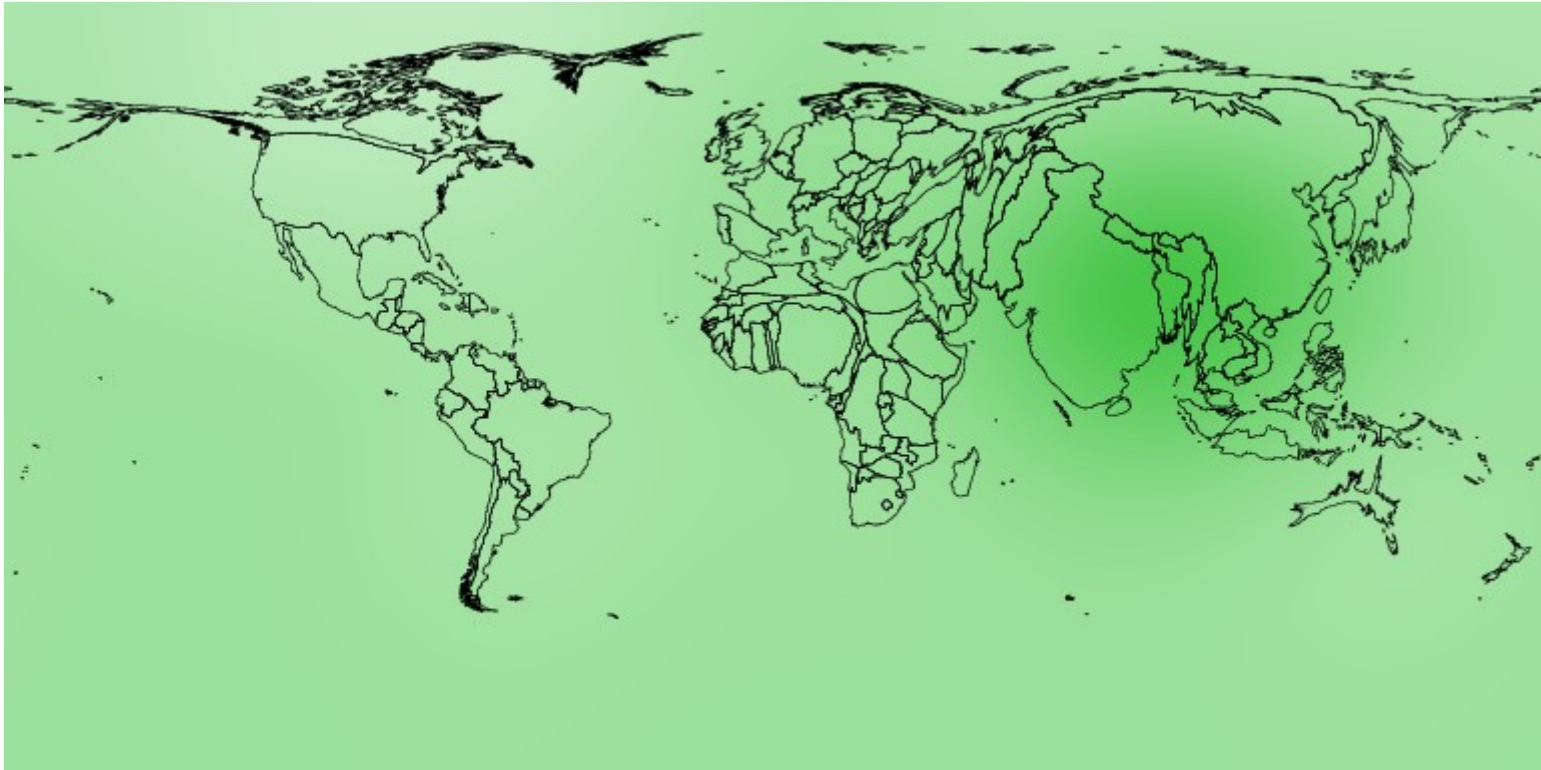














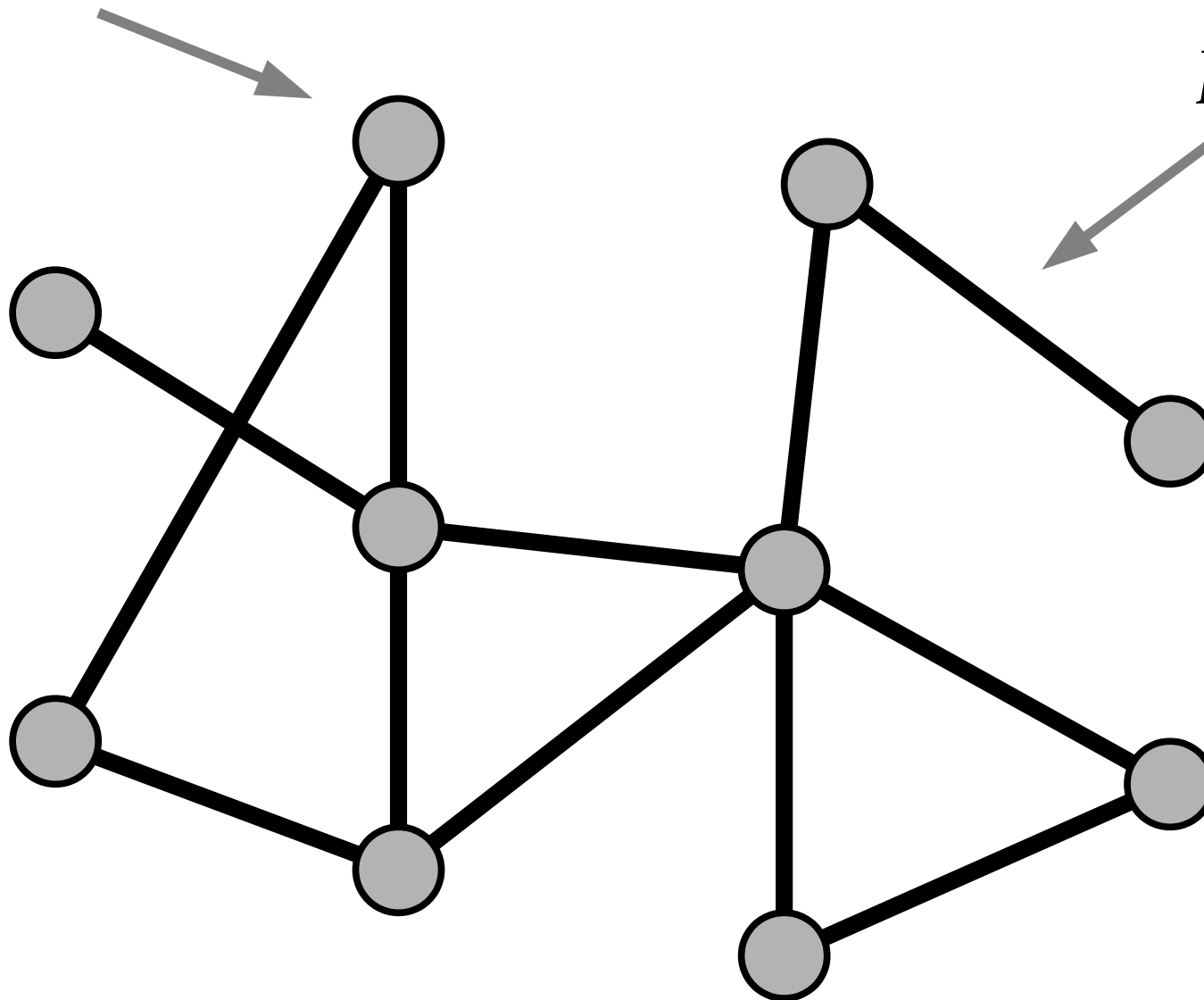


Examples

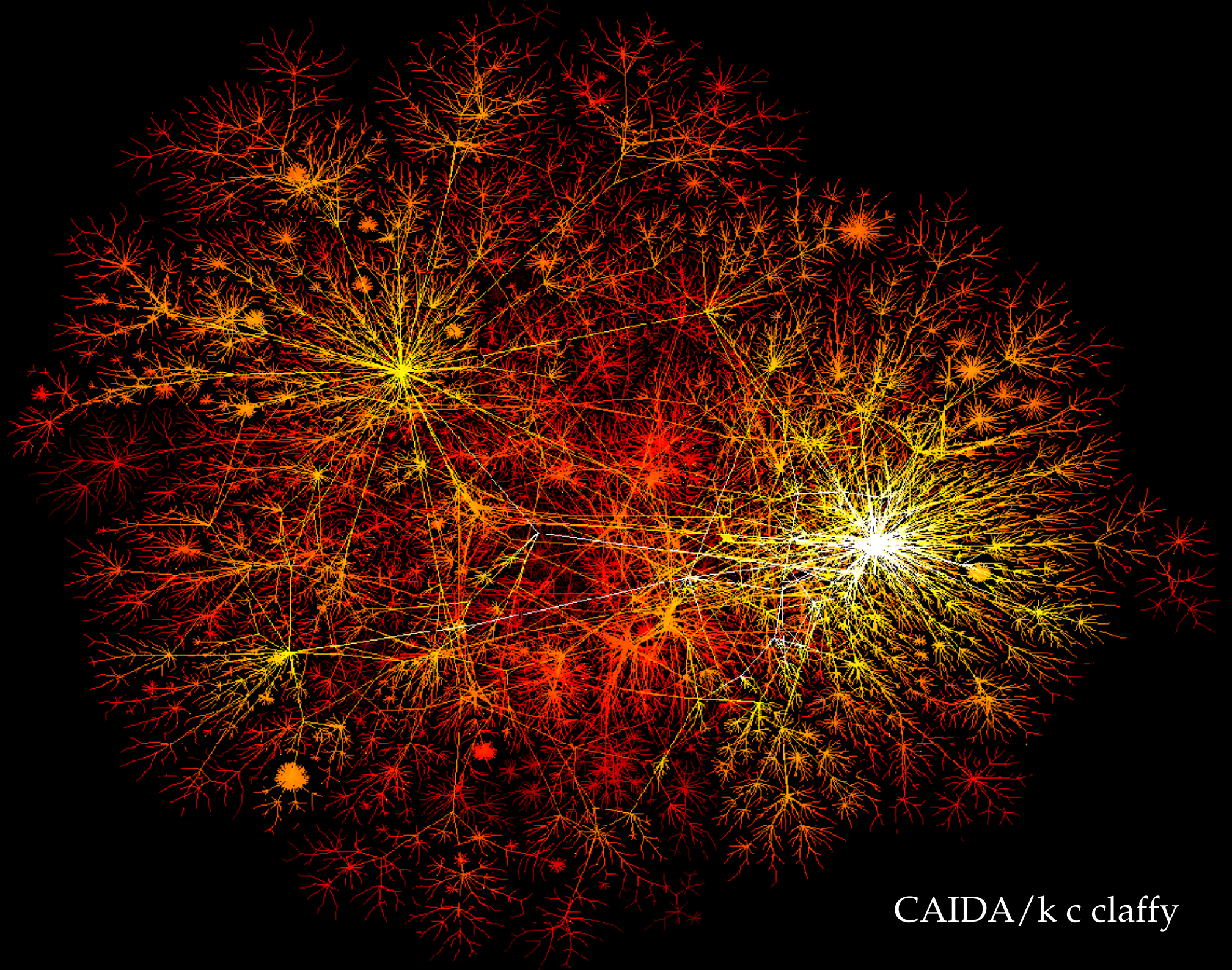
- A large selection of example maps made using this technique can be found at:

<http://www.worldmapper.org>

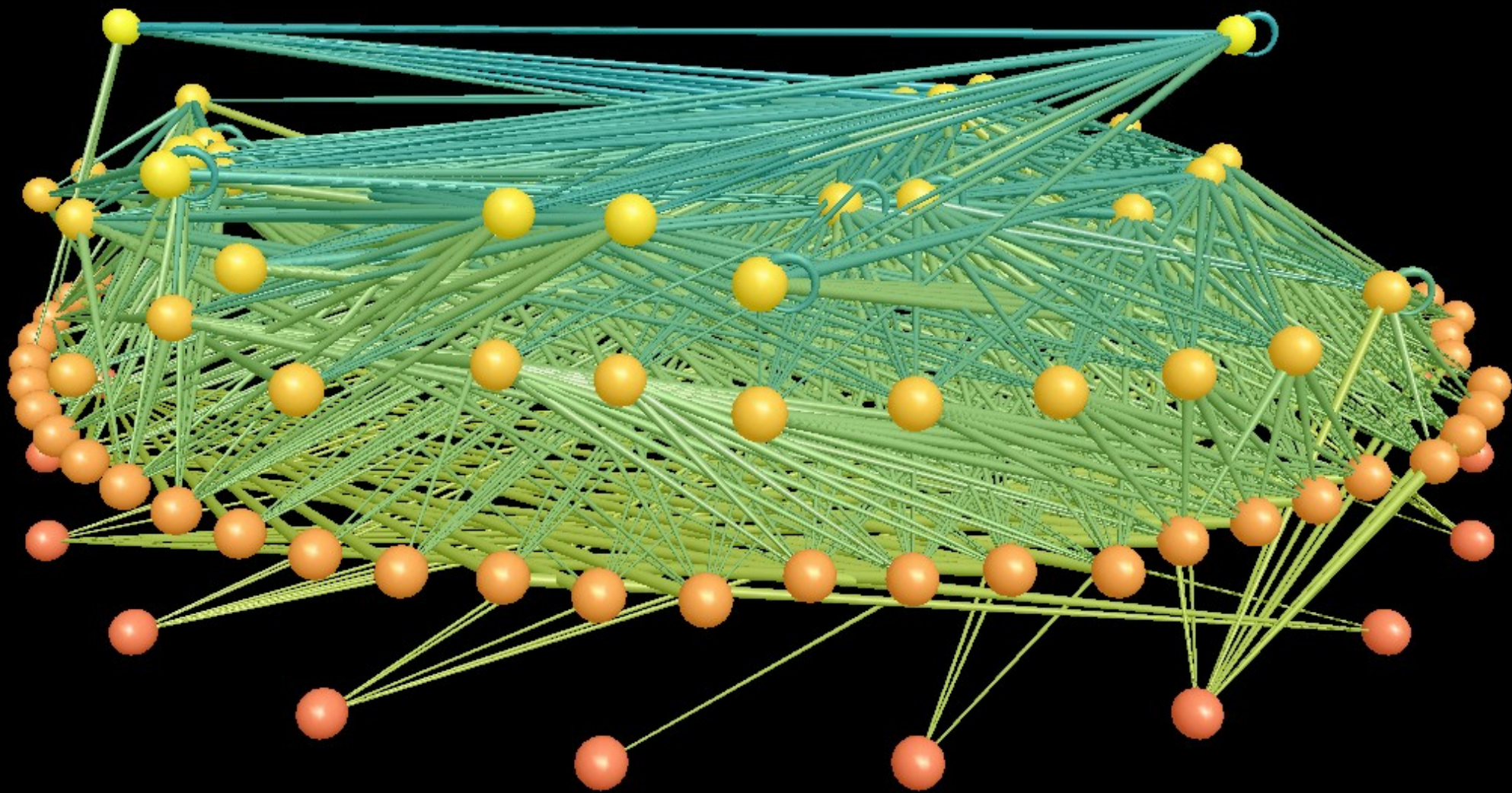
Node or vertex



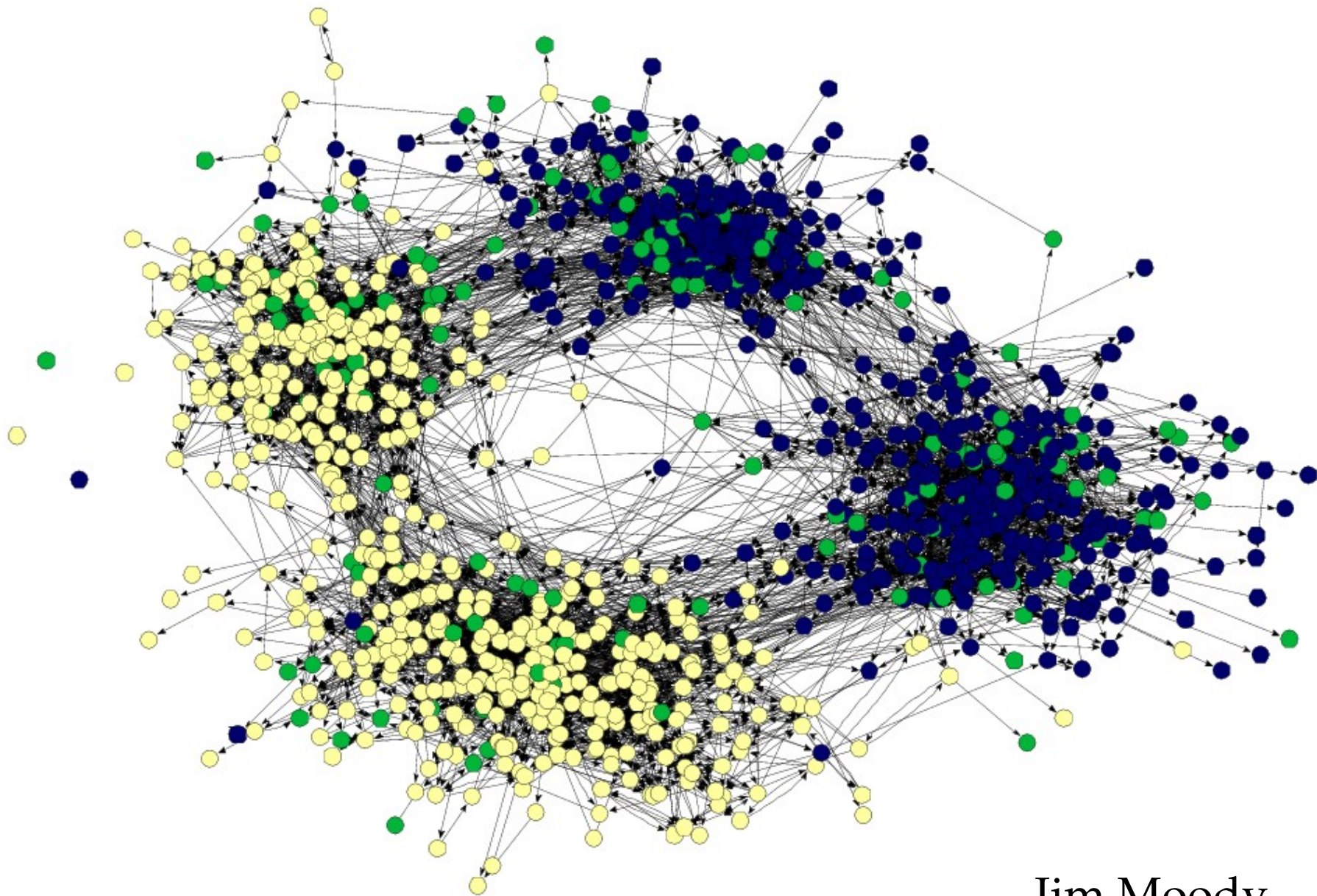
Edge



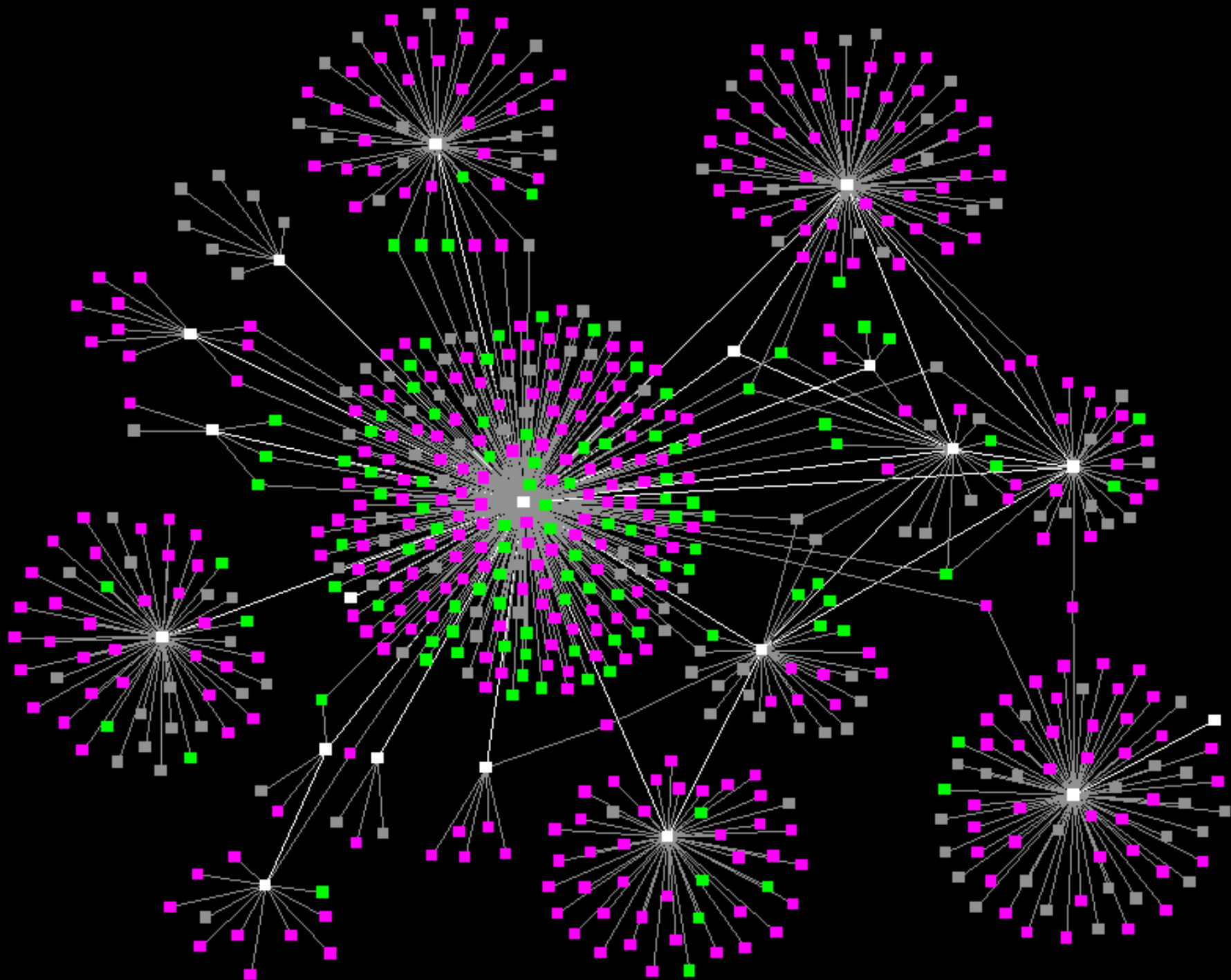
CAIDA/k c claffy



Neo Martinez and Rich Williams



Jim Moody

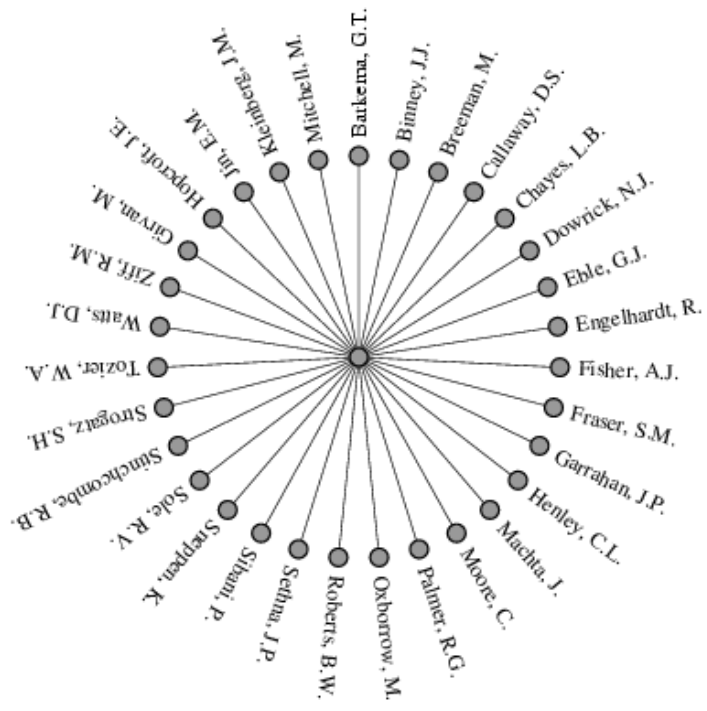


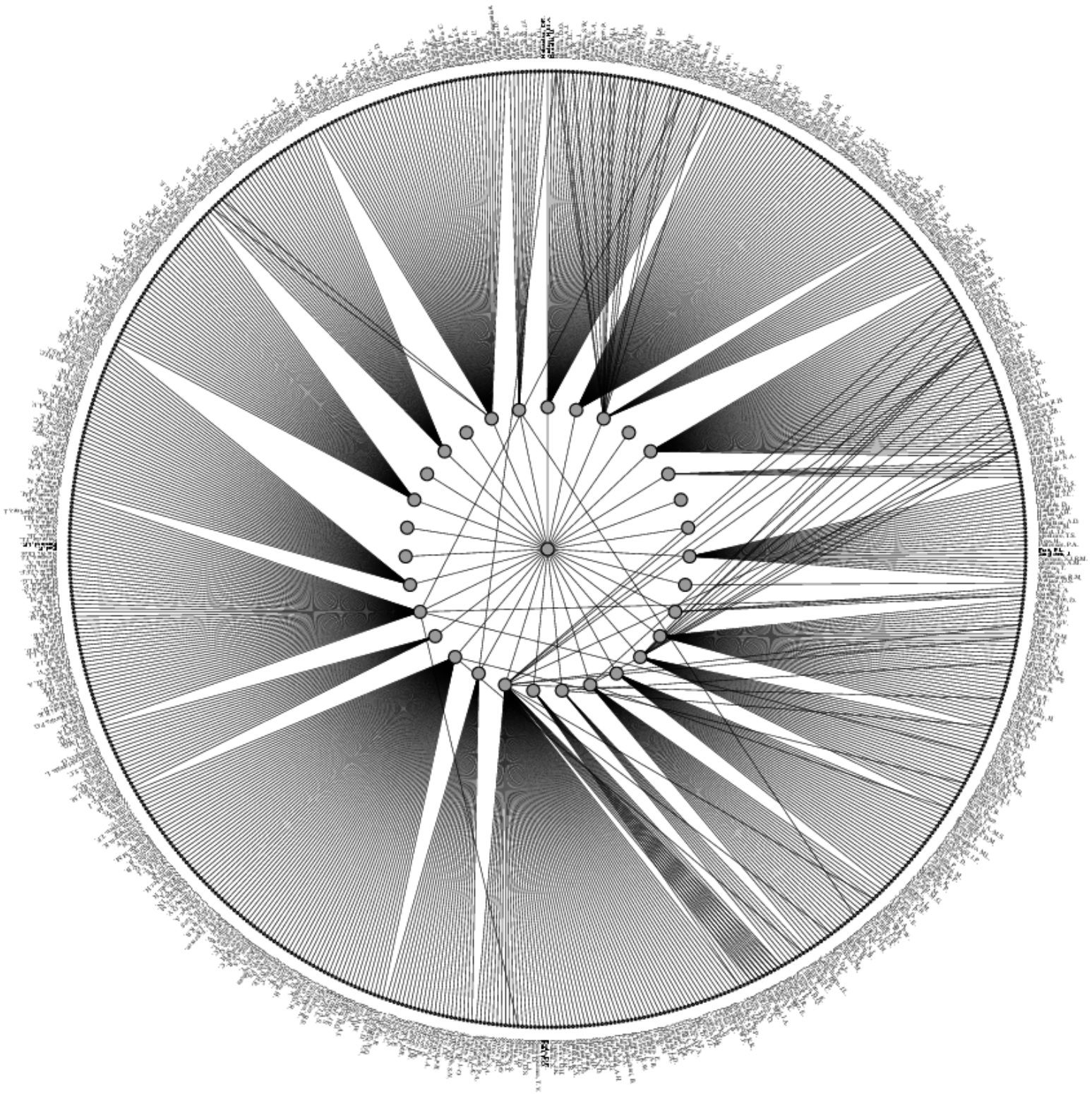
Valdis Krebs

The small-world effect

- Stanley Milgram's 1967 experiment:
 - 296 people were asked to get a letter to a target person in Boston (196 from Nebraska and 100 from Boston)
 - letters could only be passed along a chain of first-name acquaintances
 - 64 letters arrived (29%)
 - they took an average of 6.2 steps to get there
 - about a third went through one particular acquaintance of the target person (his tailor)

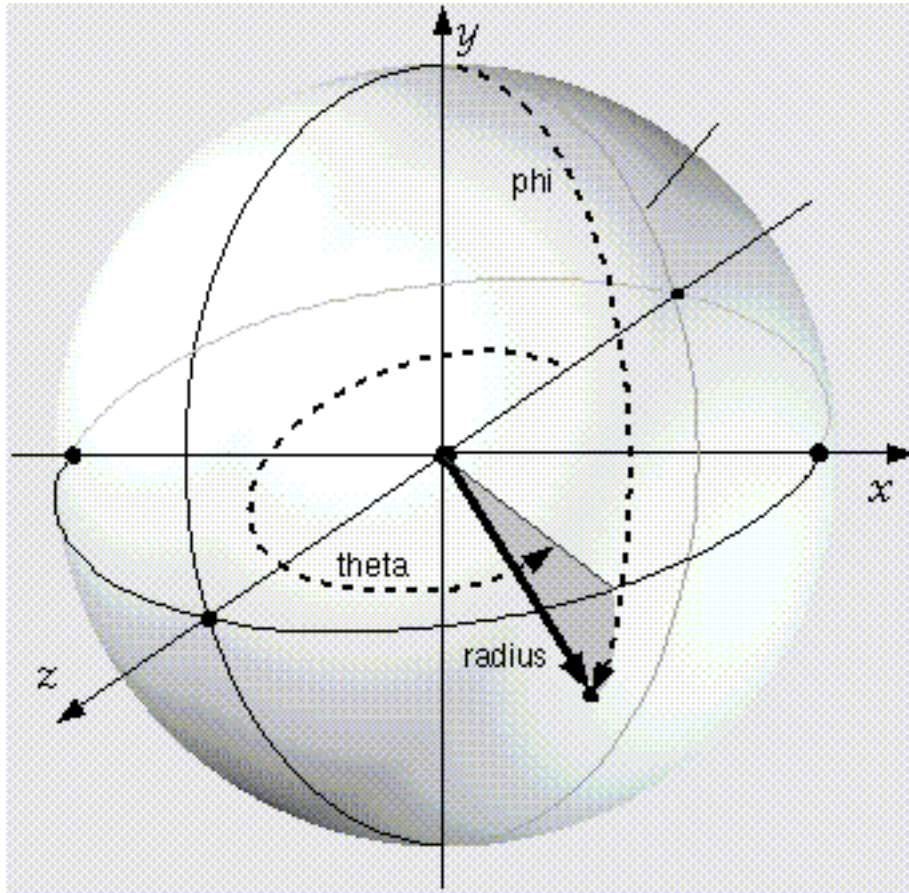
● Newman, M.E.J.





The small-world effect

- If each person knows 100 people then:
 - Number of people 1 step away from you is 100
 - People 2 steps away is $100 \times 100 = 10,000$
 - People 3 steps away is $100 \times 100 \times 100 = 1,000,000$
 - People 4 steps away is 100,000,000
 - People 5 steps away is 10,000,000,000
- But 10 billion is more than the total number of people in the world



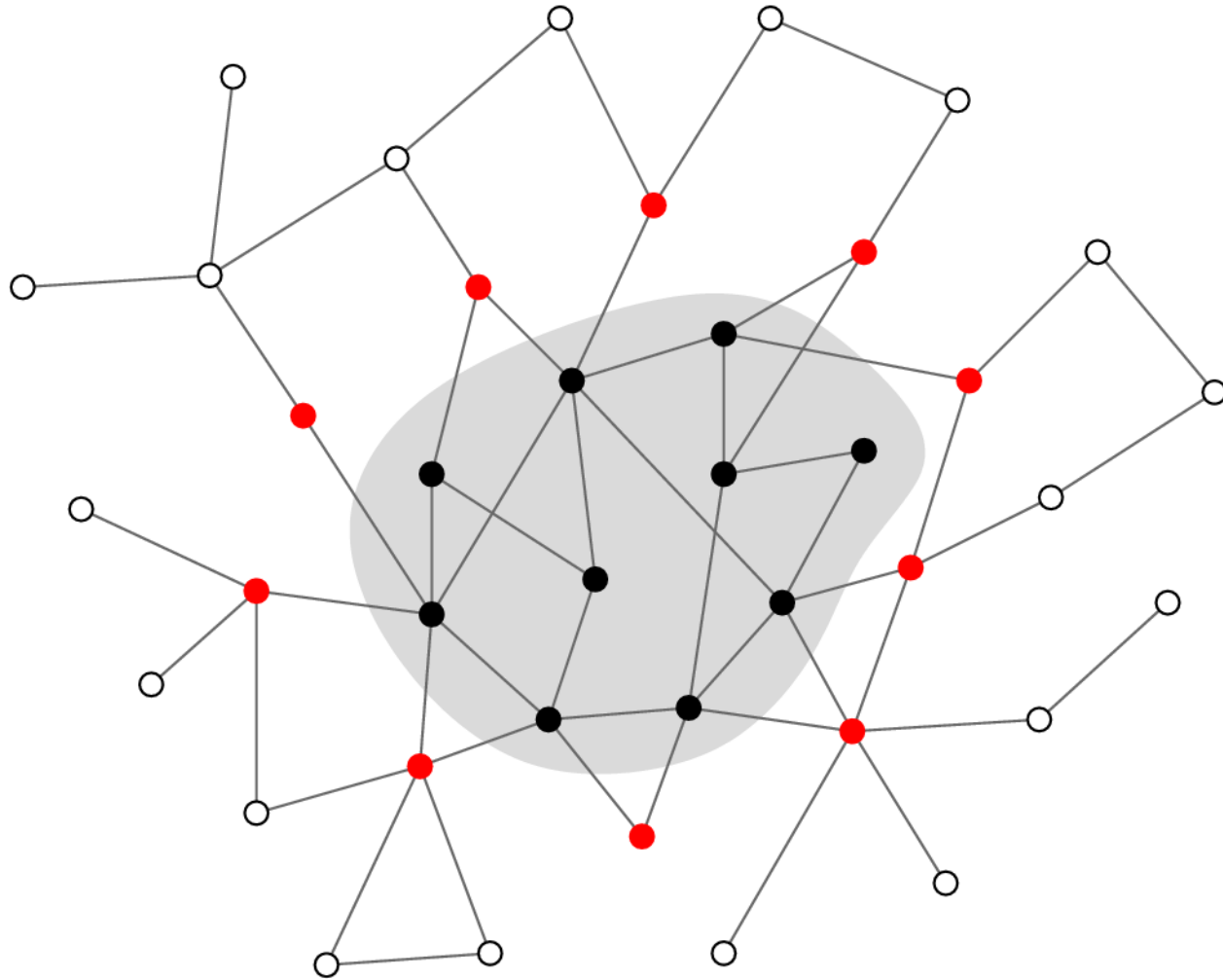
$$A = 4\pi r^2$$

$$V = \frac{4}{3}\pi r^3$$

$$A = \sqrt[3]{36\pi} V^{2/3}$$

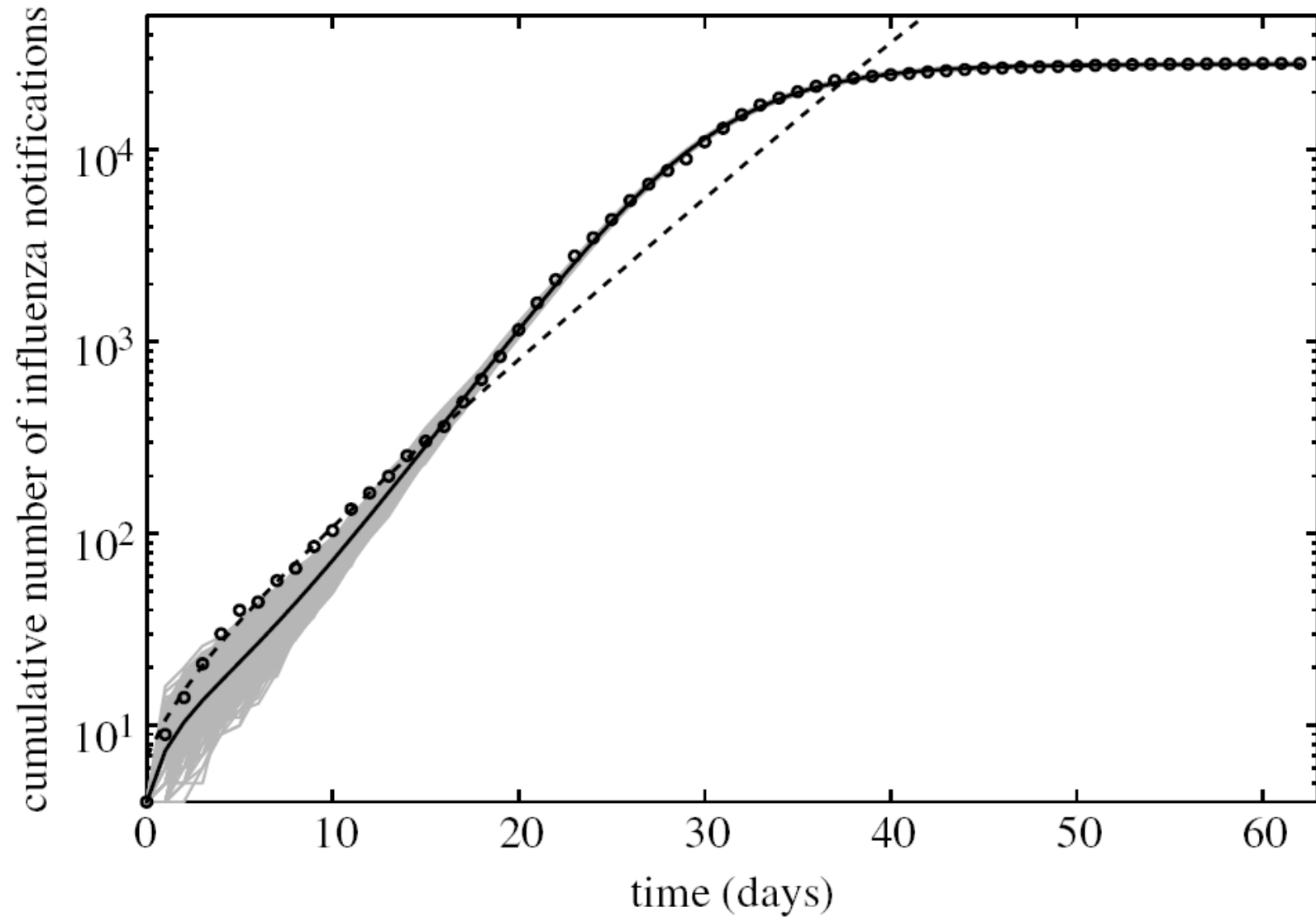
And in general dimension d

$$A \propto V^{(d-1)/d}$$

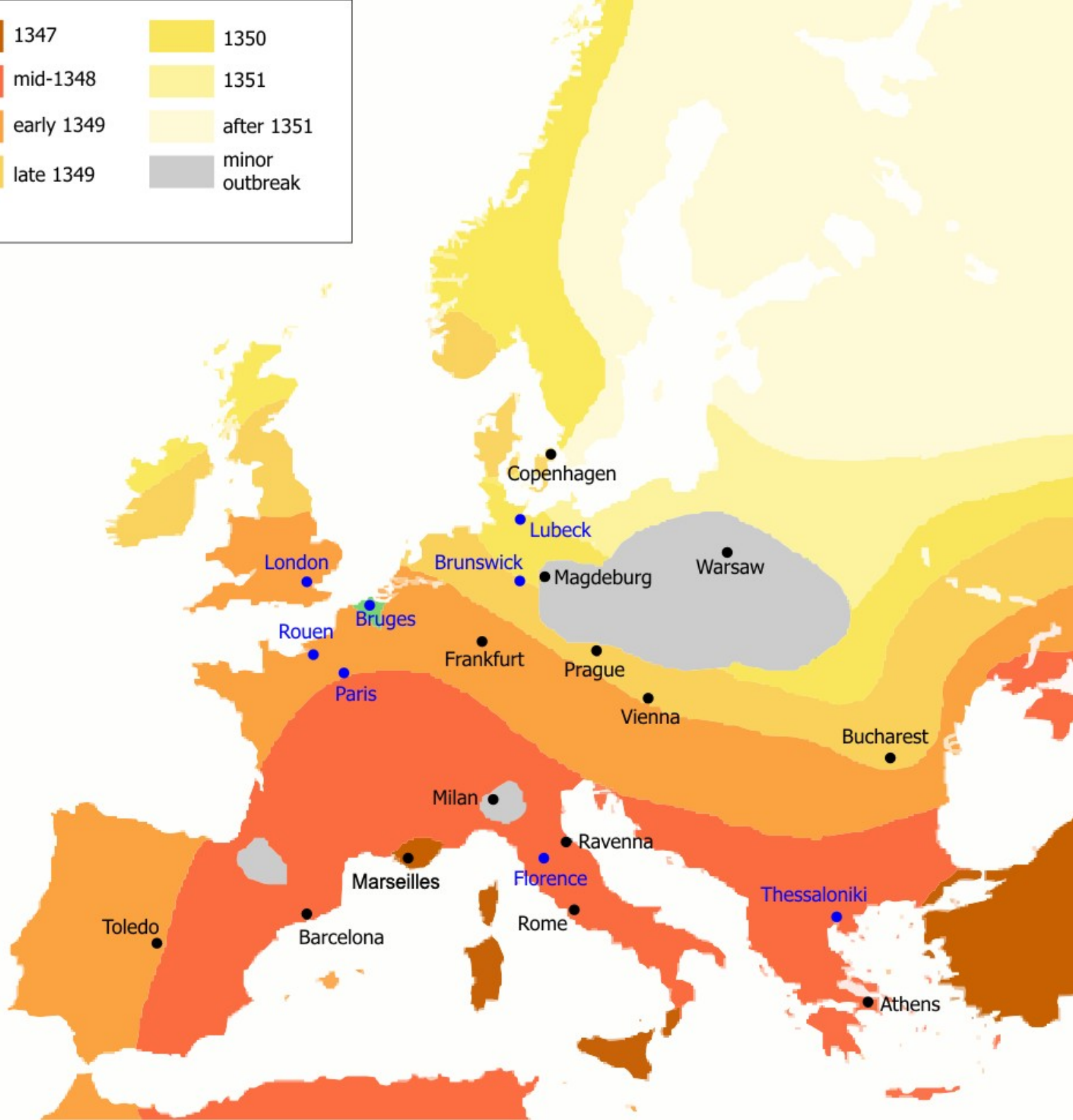
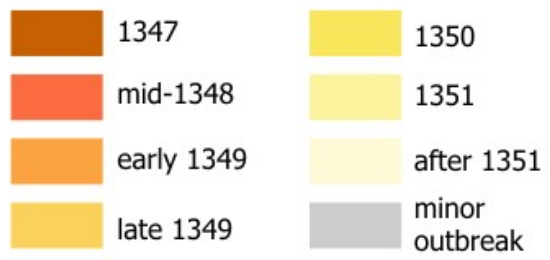


- In a network the surface can be proportional to volume
- Equivalent to choosing d infinite

Spanish flu, San Francisco 1918-1919



Chowell *et al.* 2007



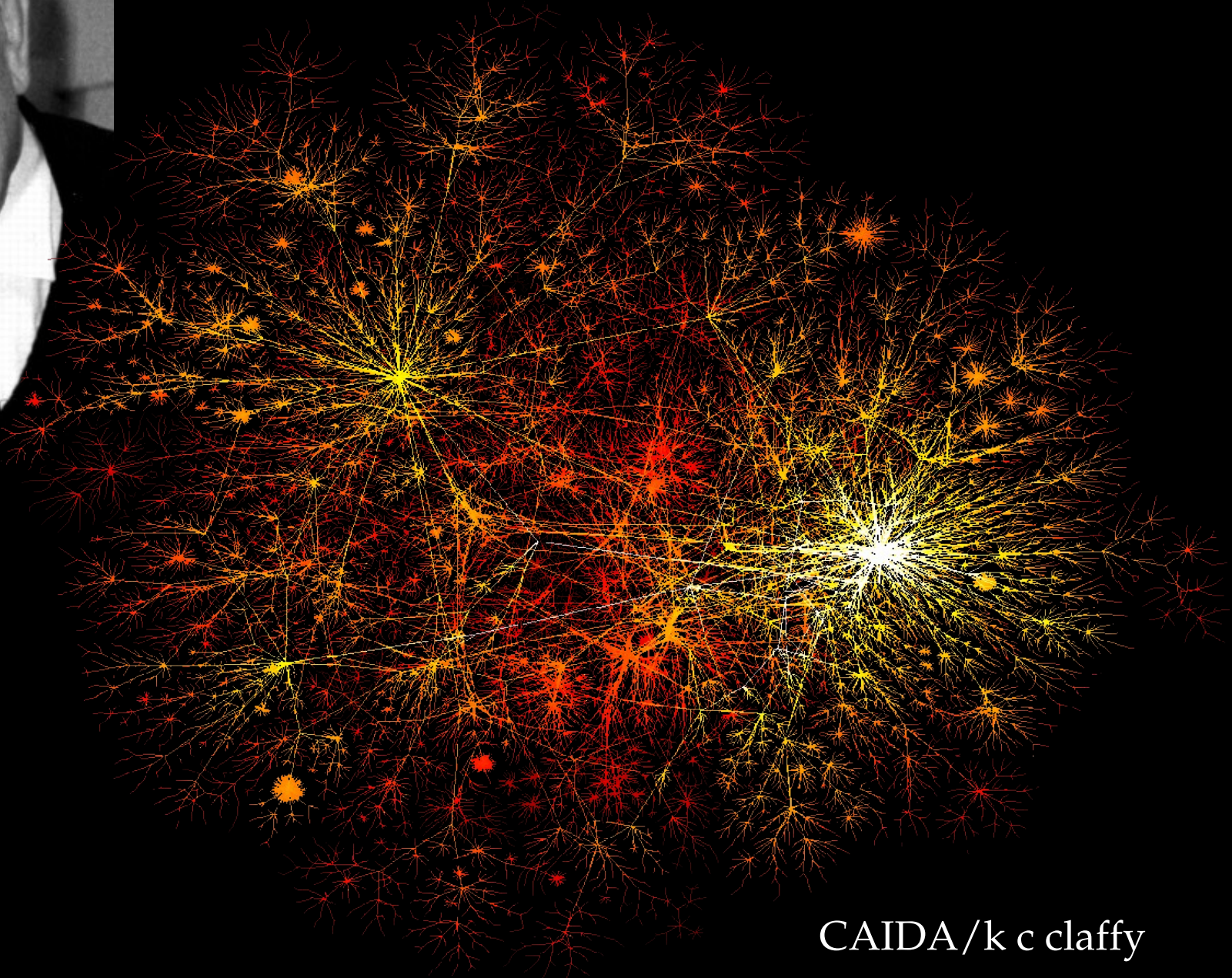
How many people does a person know?

- Just asking people is a bad way to measure this
 - People are very bad at recalling all the people they know
- Bernard and Killworth suggested a different strategy:
 - Give people a list of surnames
 - Ask them how many people they know with names on the list
 - Scale up to all the surnames in the nation (allowing for the correct [Pareto] distribution of surnames)

Top US Surnames

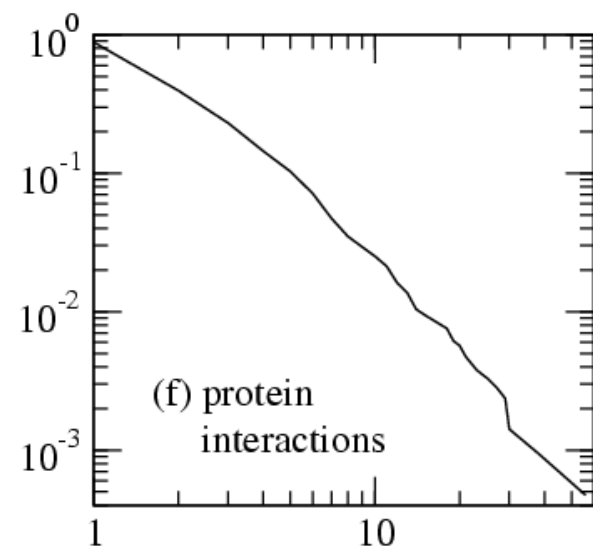
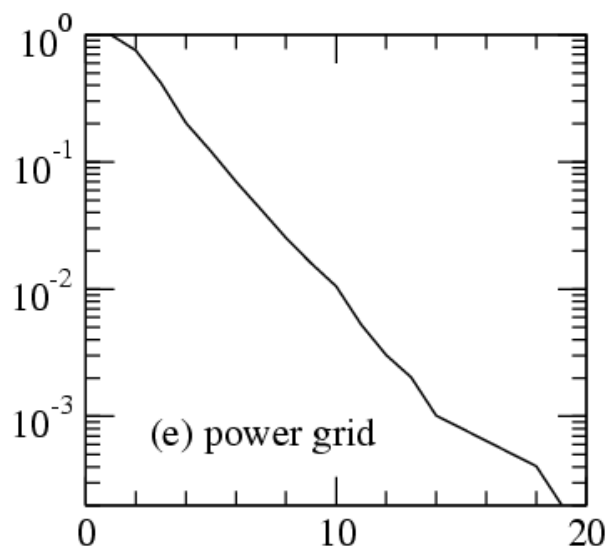
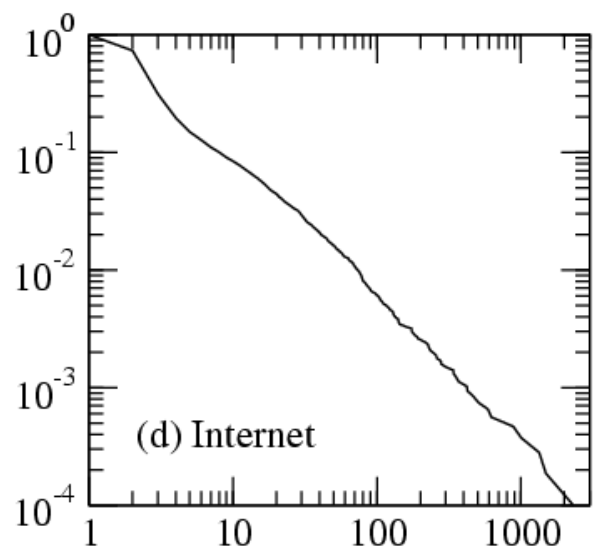
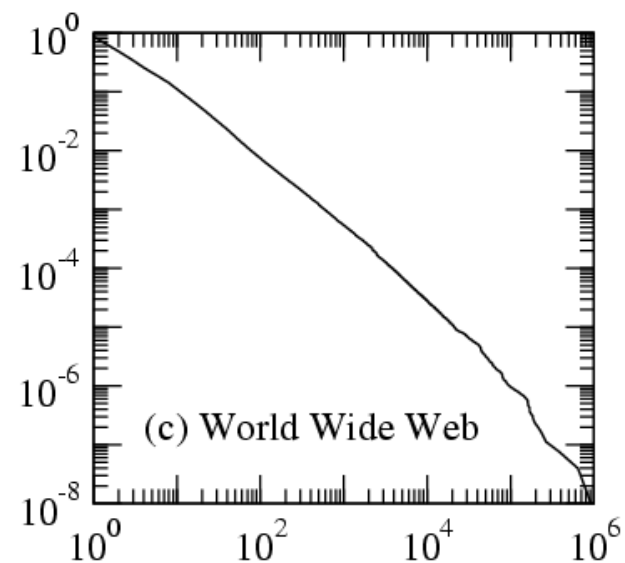
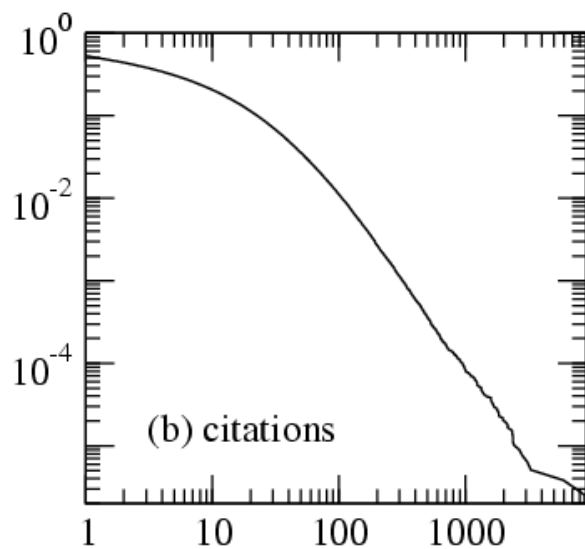
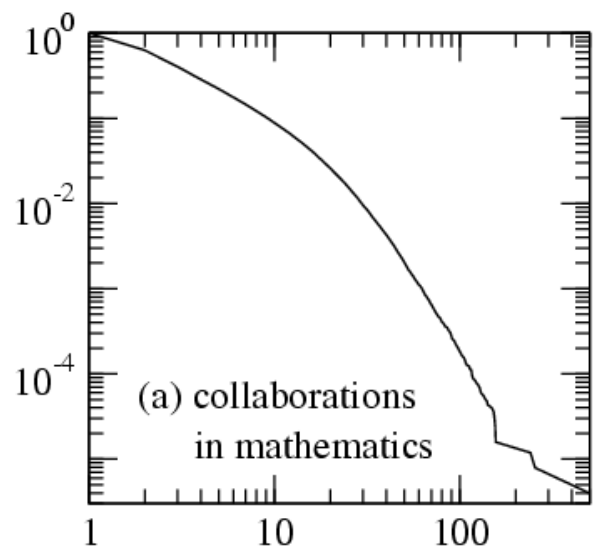
Adams	Garcia	Lewis	Robinson
Allen	Gonzalez	Lopez	Rodriguez
Anderson	Green	Martin	Scott
Baker	Hall	Martinez	Smith
Brown	Harris	Miller	Taylor
Campbell	Hernandez	Mitchell	Thomas
Carter	Hill	Moore	Thompson
Clark	Jackson	Nelson	Turner
Collins	Johnson	Parker	Walker
Davis	Jones	Perez	White
Edwards	King	Phillips	Williams
Evans	Lee	Roberts	Wilson

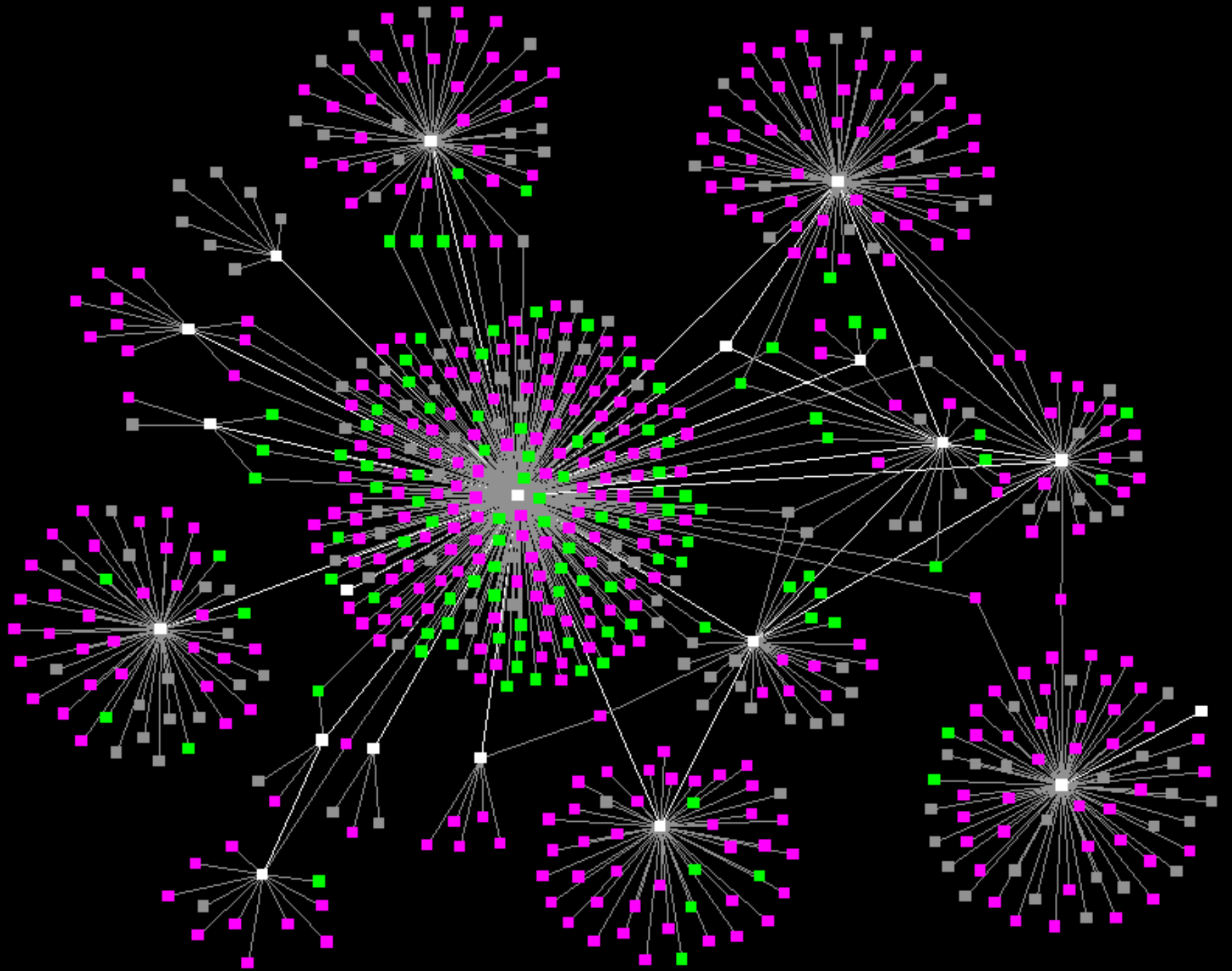
Anatol Rapoport

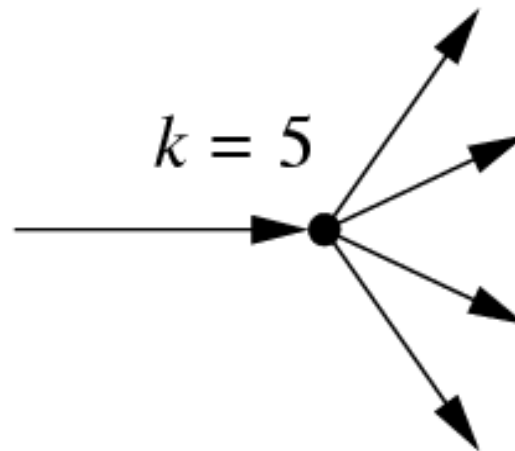


CAIDA/k c claffy

Degree distributions







Number of secondary cases is $T(k - 1)$ where T is the transmissibility.

$$R_0 = T \frac{\sum_i k_i(k_i - 1)}{\sum_i k_i} = T \frac{\langle k^2 \rangle - \langle k \rangle}{\langle k \rangle}.$$

Disease spreads if $R_0 > 1$ or if

$$T > \frac{\langle k \rangle}{\langle k^2 \rangle - \langle k \rangle}.$$

Epidemic threshold depends on the structure of the network.

Information content of radio transmissions

Suppose we send a message composed of a stream of symbols or characters, and suppose that the character A occurs with probability p_A .

- If $p_A = \frac{1}{2}$, then we need 1 bit to say whether it's present.
- If $p_A = \frac{1}{4}$, then we need 2 bits.
- If $p_A = \frac{1}{8}$, then we need 3 bits.

In general a character i that occurs with probability p_i carries information $\log_2(1/p_i) = -\log_2 p_i$ bits.

Average information per character

$$S = - \sum_i p_i \log_2 p_i.$$

(Shannon 1948.)

What distribution over characters gives the highest mean rate of transmission of information? We maximize the average information S subject to the constraint

$$\sum_i p_i = 1,$$

giving

$$-\frac{\partial}{\partial p_i} \sum_i p_i \log p_i - \alpha \frac{\partial}{\partial p_i} \left[\sum_i p_i - 1 \right] = 0,$$

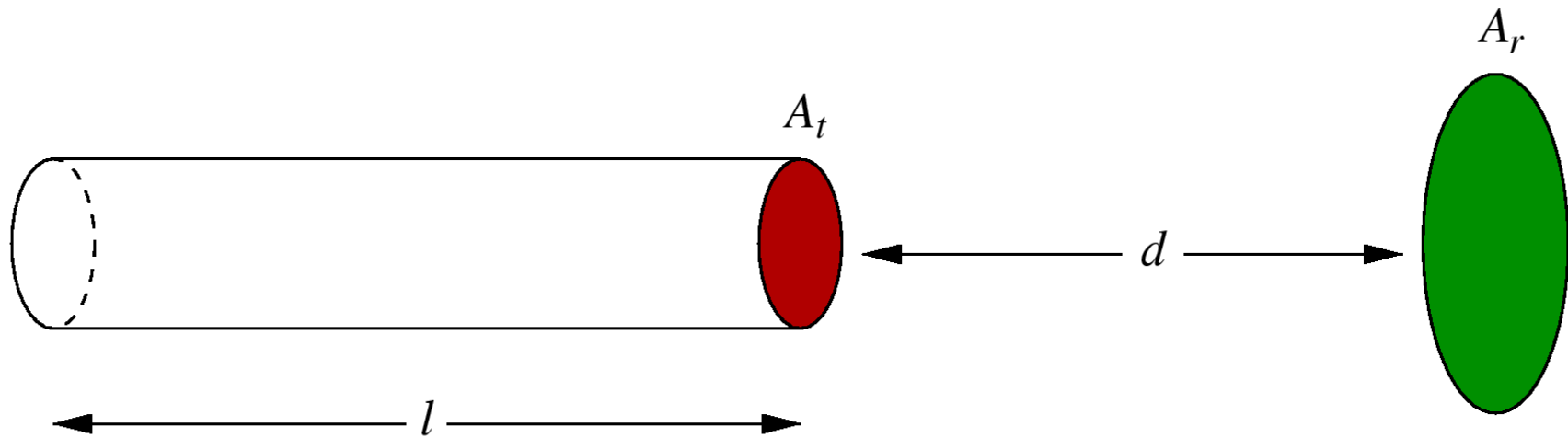
for all i , or

$$p_i = e^{-\alpha-1} = \text{constant}.$$

So all characters are equally likely.

A message with maximal information rate is indistinguishable from random noise to an observer unfamiliar with the code in which the message is written.

To study the equivalent question for radio waves, consider the following thought experiment:



We set up electromagnetic microstates in the perfectly reflecting interior of the cavity. Each microstate is one “character” of the message.

We have a mean transmission power or “energy budget” of $\langle E \rangle$ per message.

Any number of photons is allowed in the cavity.

Constraints:

$$\sum_i p_i = 1, \quad \sum_i p_i E_i = \langle E \rangle, \quad \sum_i p_i N_i = \langle N \rangle.$$

This is a well-known problem:

$$\frac{\partial S}{\partial p_i} - \alpha \frac{\partial}{\partial p_i} \left[\sum_i p_i - 1 \right] - \beta \frac{\partial}{\partial p_i} \left[\sum_i p_i E_i - \langle E \rangle \right] - \gamma \frac{\partial}{\partial p_i} \left[\sum_i p_i N_i - \langle N \rangle \right] = 0$$

for all i .

$$p_i = \frac{e^{-\beta(E_i - \mu N_i)}}{Z},$$

where $\mu = \gamma/\beta$ and $Z = \sum_i e^{-\beta(E_i - \mu N_i)}$.

The rest of the calculation is well known. We put

$$N = \sum_k n_k, \quad E = \sum_k n_k \epsilon_k,$$

and perform the sum over all the occupancy numbers n_k to find that

$$Z = \sum_{\{n_k\}} \prod_k e^{-\beta(\epsilon_k - \mu)n_k} = \prod_k \sum_{n_k=0}^{\infty} e^{-\beta(\epsilon_k - \mu)n_k} = \prod_k [1 - e^{-\beta(\epsilon_k - \mu)}]^{-1}.$$

Then the mean number of photons in the cavity is

$$\begin{aligned} \langle N \rangle &= \beta^{-1} \frac{\partial \ln Z}{\partial \mu} = \sum_k \frac{1}{e^{\beta(\epsilon_k - \mu)} - 1} \\ &= \sum_k \langle n_k \rangle. \end{aligned}$$

So

$$\langle n_k \rangle = \frac{1}{e^{\beta\epsilon_k} - 1}.$$

This distribution over modes maximizes the information content.

But now here's the trick. Instead of passing the cavity to the receiver, we just take the lid off and send the message by radio.

The density of single-particle states that have momentum in the right direction to hit the receiver distance d away is

$$\rho(\epsilon) = \frac{2\ell A_t A_r \epsilon^2}{d^2 h^3 c^3}.$$

So the power spectrum of the message is

$$I(\epsilon) = \frac{2\ell A_t A_r}{d^2 h^3 c^3} \frac{\epsilon^3}{e^{\beta\epsilon} - 1}.$$

*A message with maximal information rate transmitted using radio waves **has the energy spectrum of blackbody radiation.***

Apparent temperature:

$$T^4 = \frac{15h^3c^2}{2\pi^4} \frac{d^2}{A_t A_r} P.$$

Information rate:

$$\frac{dS}{dt} = \frac{8\pi^4}{45h^3c^2} \frac{A_t A_r}{d^2} T^3 = \left[\frac{512\pi^4}{1215h^3c^2} \frac{A_t A_r}{d^2} P^3 \right]^{\frac{1}{4}}.$$

For a transmitter and receiver of one square meter each, a meter apart, with a power of $P = 1$ W, the information rate is 1.61×10^{21} bits per second.

Note that the information rate increases with the area of the transmitter and receiver, so that the best information rate for a given energy budget is achieved for large antennas and low apparent temperature.



Any sufficiently advanced technology is
indistinguishable from noise

